



# NATIONAL CYBER SUMMIT

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# U Mad? Binary Analysis with the Angr Framework

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# Intro

- What is binary analysis?
- What is angr? (An unapologetic oversimplification)
- Demos!





# What is Binary Analysis?

- Software bugs have taken down spaceships<sup>1</sup>, caused nuclear centrifuges to spin out of control<sup>2</sup>, and forced the recall of 100,000s of vehicles resulting in billions of dollars in damages<sup>3</sup>.
- How can you find these bugs when source code is unavailable?
- **Reverse Engineering, Vulnerability Assessment, and Binary Analysis**
- Process: Disassemble, Triage, Understand, Analyze, Symbolize.

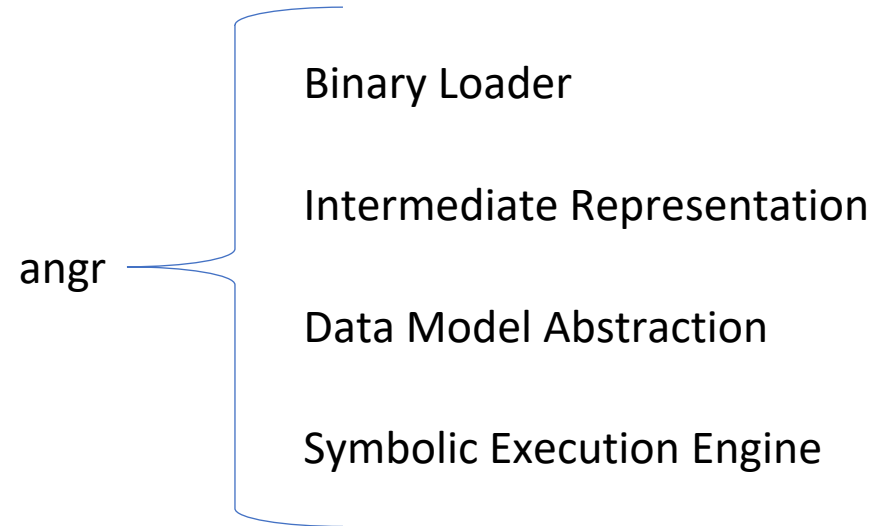
<sup>1</sup> Ariane 5: Who Dunit? <https://ieeexplore.ieee.org/document/589224/>

<sup>2</sup> Lessons from Stuxnet <https://ieeexplore.ieee.org/document/5742014/>

<sup>3</sup> A Case Study of Toyota Unintended Acceleration and Software Safety [https://users.ece.cmu.edu/~koopman/pubs/koopman14\\_toyota\\_ua\\_slides.pdf](https://users.ece.cmu.edu/~koopman/pubs/koopman14_toyota_ua_slides.pdf)

# What is angr?

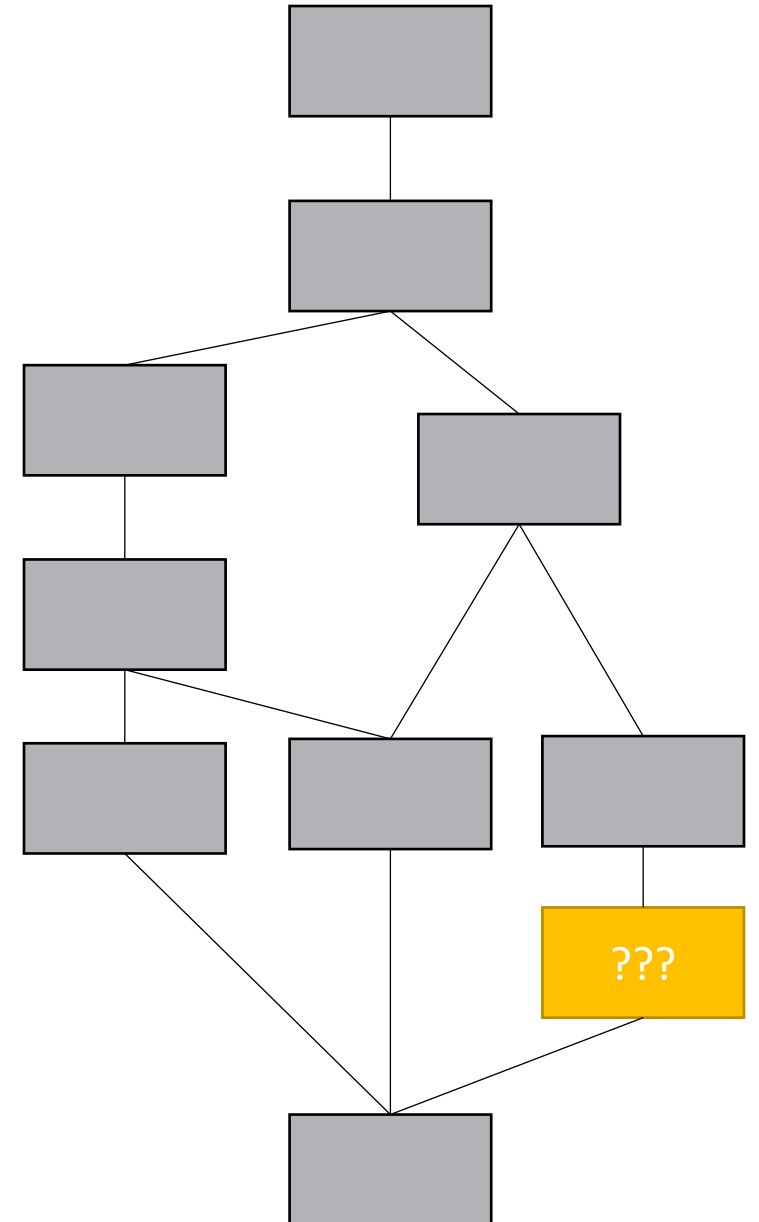
- Binary Analysis Framework developed by the University of California Santa Barbara since 2013.
- Features:
  - iPython accessible
  - Powerful analyses
  - Versatile
  - Open and expandable
  - Architecture “independent”



# Software Analysis

“How do I trigger path  $x$  or condition  $y$ ?”

- Dynamic analysis
  - Input  $a$ ? No. Input  $b$ ? No. Input  $c$ ? ...
  - Based on concrete inputs to application
- Static analysis
  - “You can’t”
  - “You might be able to...”
  - “IDK”
  - Based on various static techniques.



# Symbolic Execution

“How do I trigger path  $x$  or condition  $y$ ?”

1. Interpret the application.
2. Track “constraints” on variables.
3. When the required condition is triggered, “concretize” to obtain a possible input.



# Symbolic Execution Example

```
def f (x, y):  
    if (x > y):  
        x = x + y  
        y = x - y  
        x = x - y  
        if (x - y > 0):  
            call g()  
    return (x, y)
```

This function swaps the values of x and y when  $x > y$ .

The  $x - y > 0$  statement is always false so the call is unreachable.

Source code is here but our techniques allow for the same analysis without source code.



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Execute the program on *symbolic values*.

# Symbolic Execution Example

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*Symbolic state* maps variables to symbolic values.

$x \rightarrow A$

$y \rightarrow B$

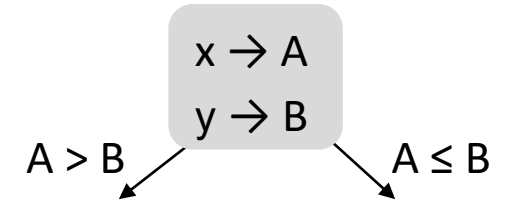
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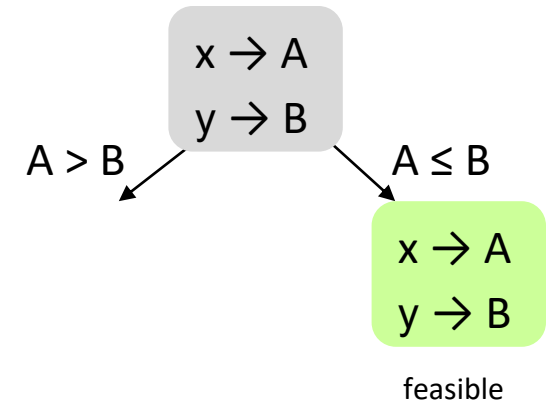
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All paths in the program form its *execution tree*, in which some paths are *feasible* and some are *infeasible*.



# Symbolic Execution Example

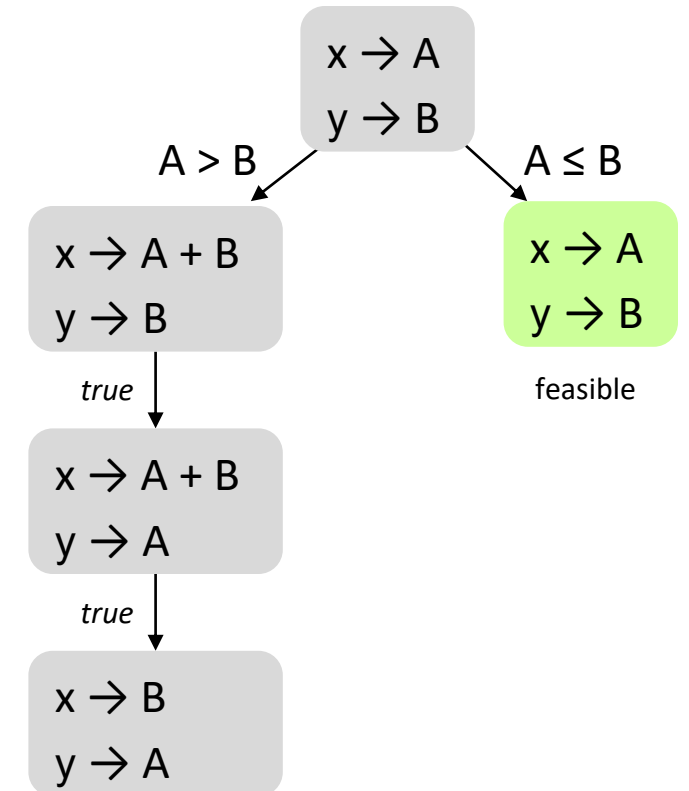
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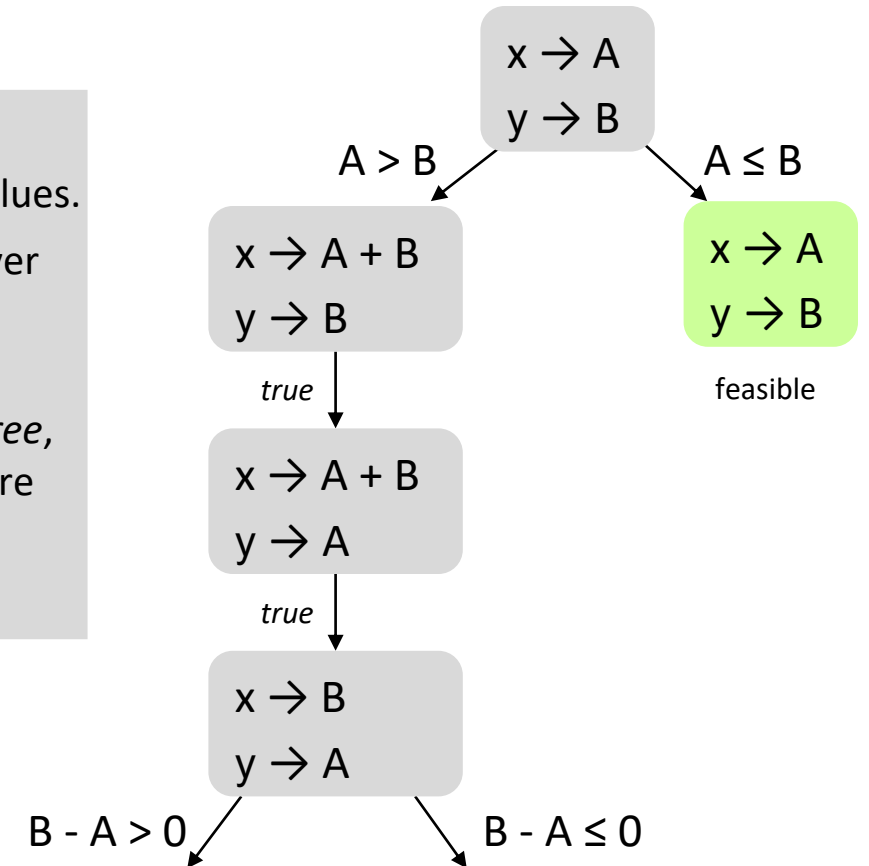
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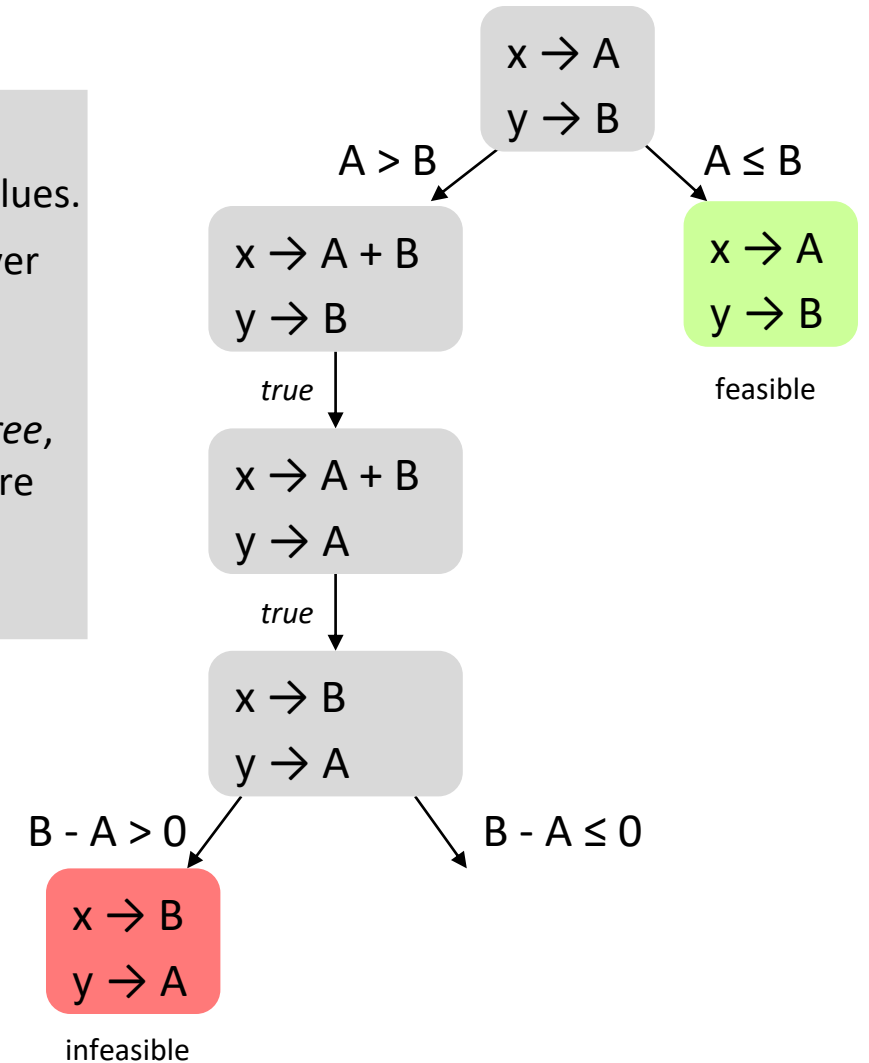
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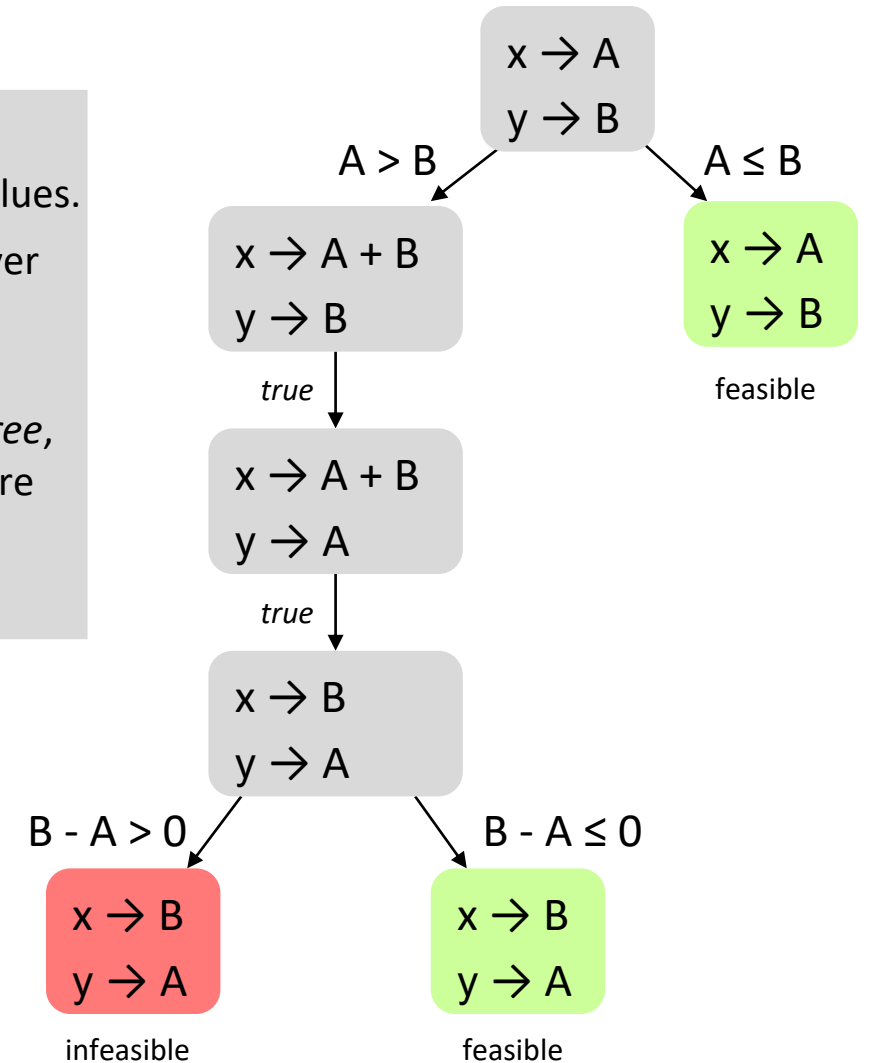
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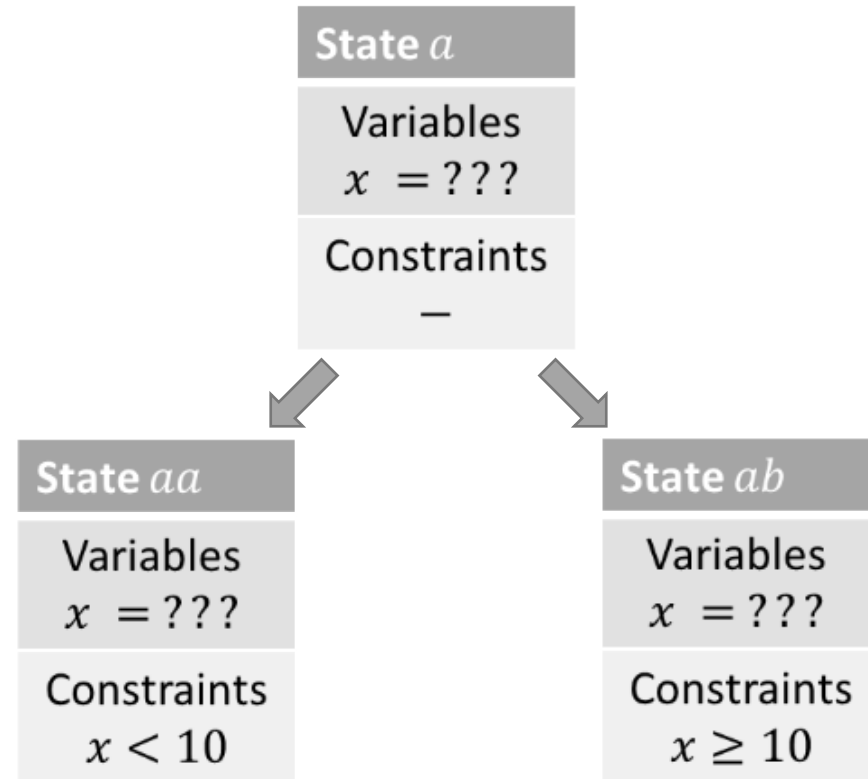
# Symbolic Execution Example

```
x = int (input())
if x ≤ 10:
    if x < 100:
        print "Two!"
    else:
        print "Lots!"
else:
    print "One!"
```

State <i>a</i>
Variables $x = ???$
Constraints —

# Symbolic Execution Example

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x = int (input())  
if x ≤ 10:  
    if x < 100:  
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# Symbolic Execution Example

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x = int (input())
if x ≤ 10:
    if x < 100:
        print "Two!"
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else:
    print "One!"
```

State *aa*

Variables

$x = ???$

Constraints

$x < 10$

State *ab*

Variables

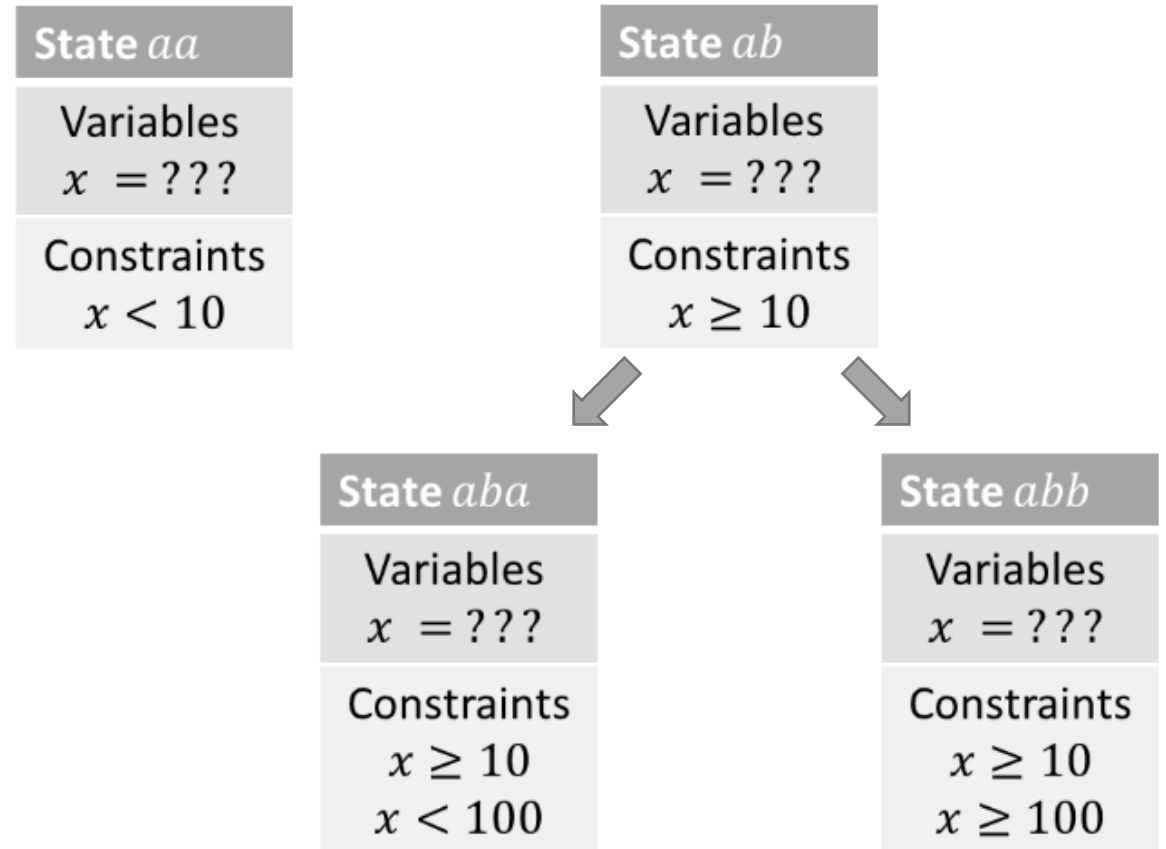
$x = ???$

Constraints

$x \geq 10$

# Symbolic Execution Example

```
x = int (input())  
if x ≤ 10:  
    if x < 100:  
        print "Two!"  
    else:  
        print "Lots!"  
else:  
    print "One!"
```





# Symbolic Execution Example

```
x = int (input())
if x ≤ 10:
    if x < 100:
        print "Two!"
    else:
        print "Lots!"
else:
    print "One!"
```

State *aba*

Variables

$x = ???$

Constraints

$x \geq 10$

$x < 100$

# Symbolic Execution Example

```
x = int (input())
if x ≤ 10:
    if x < 100:
        print "Two!"
    else:
        print "Lots!"
else:
    print "One!"
```

State *aba*

Variables

$x = ???$

Constraints

$x \geq 10$

$x < 100$



Concretized *aba*

Variables

$x = 99$

# Demo: crackme

- Available at [https://github.com/bendenton/2018\\_NCS](https://github.com/bendenton/2018_NCS)

# Demo: crackme2

- Available at [https://github.com/bendenton/2018\\_NCS](https://github.com/bendenton/2018_NCS)

# Demo: crackme3

- Available at  
[https://github.com/bendenton/2018\\_NCS](https://github.com/bendenton/2018_NCS)



# Questions?

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[https://github.com/bendenton/2018\\_NCS](https://github.com/bendenton/2018_NCS)

<https://www.linkedin.com/in/ben-denton/>

